

# What Can We Do with a Quantum Computer?

Matthias Troyer – Station Q, ETH Zurich

# Classical computers have come a long way



**Antikythera mechanism**  
astronomical positions  
(100 BC)



**Kelvin's harmonic analyzer**  
prediction of tides  
(1878)



**Difference Engine**  
(1822)



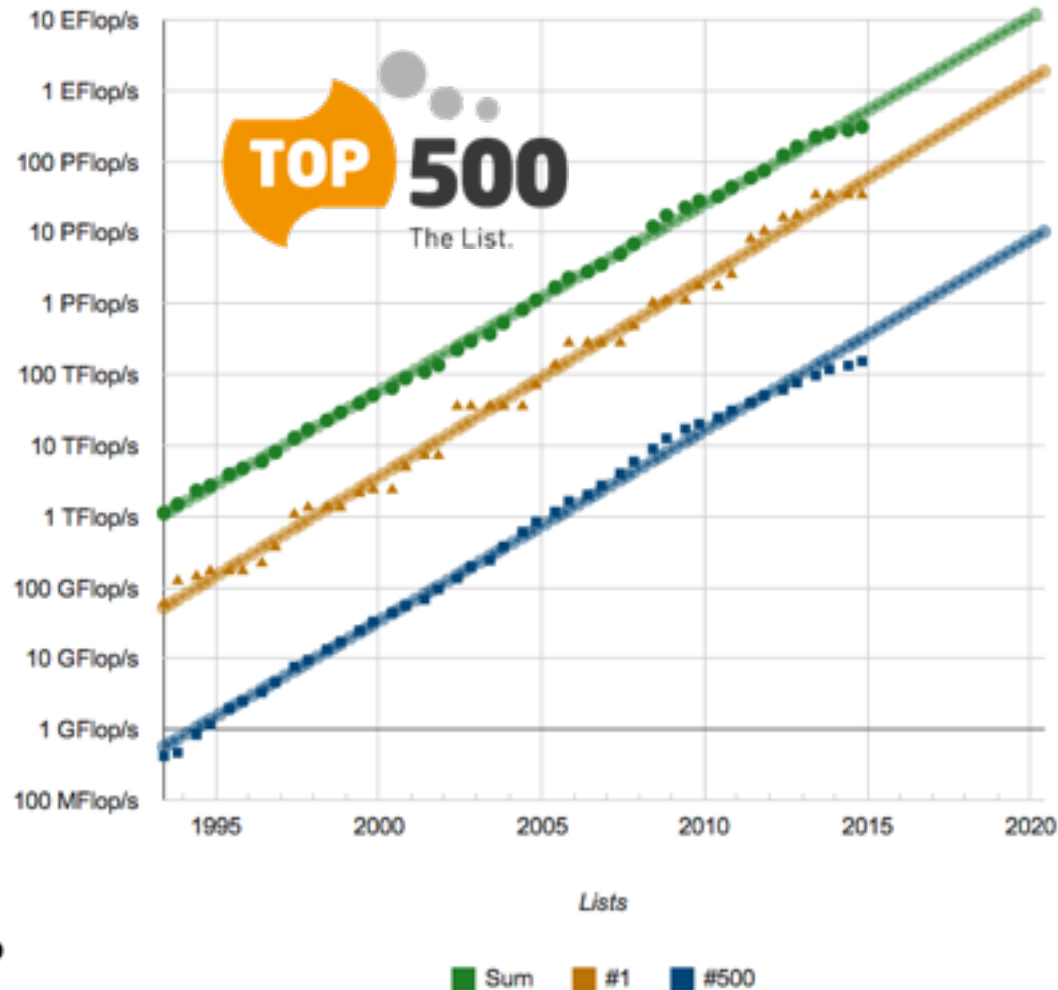
**ENIAC**  
(1946)



**Titan, ORNL**  
(2013)

Is there anything that we cannot solve on future supercomputers?

# How long will Moore's law continue?



Do we see signs of the end of Moore's law?

Can we go below 7nm feature size?

Can we use more than 3 million cores?

Can we fight the recent exponential increase in power consumption?



## Enabling technologies for beyond exascale computing

- We are not referring to  $10^{21}$  flops
- “Beyond exascale” systems as we are defining them will be based on new technologies that will finally result in the much anticipated (but unknown) phase change to truly new paradigms/methodologies.



**Paul Messina**

*Director of Science*

*Argonne Leadership Computing Facility*

*Argonne National Laboratory*

*July 9, 2014*

*Cetraro*

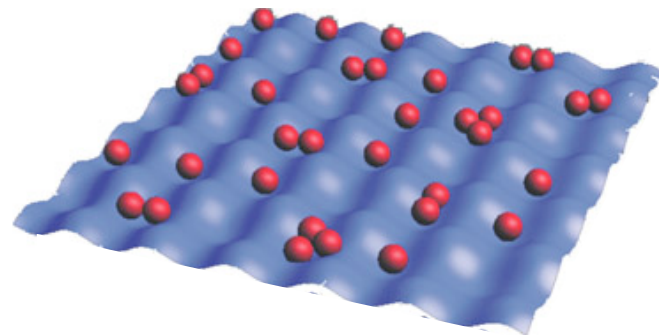
# Our bet: quantum devices



Quantum randomness



Quantum communication



Quantum simulation



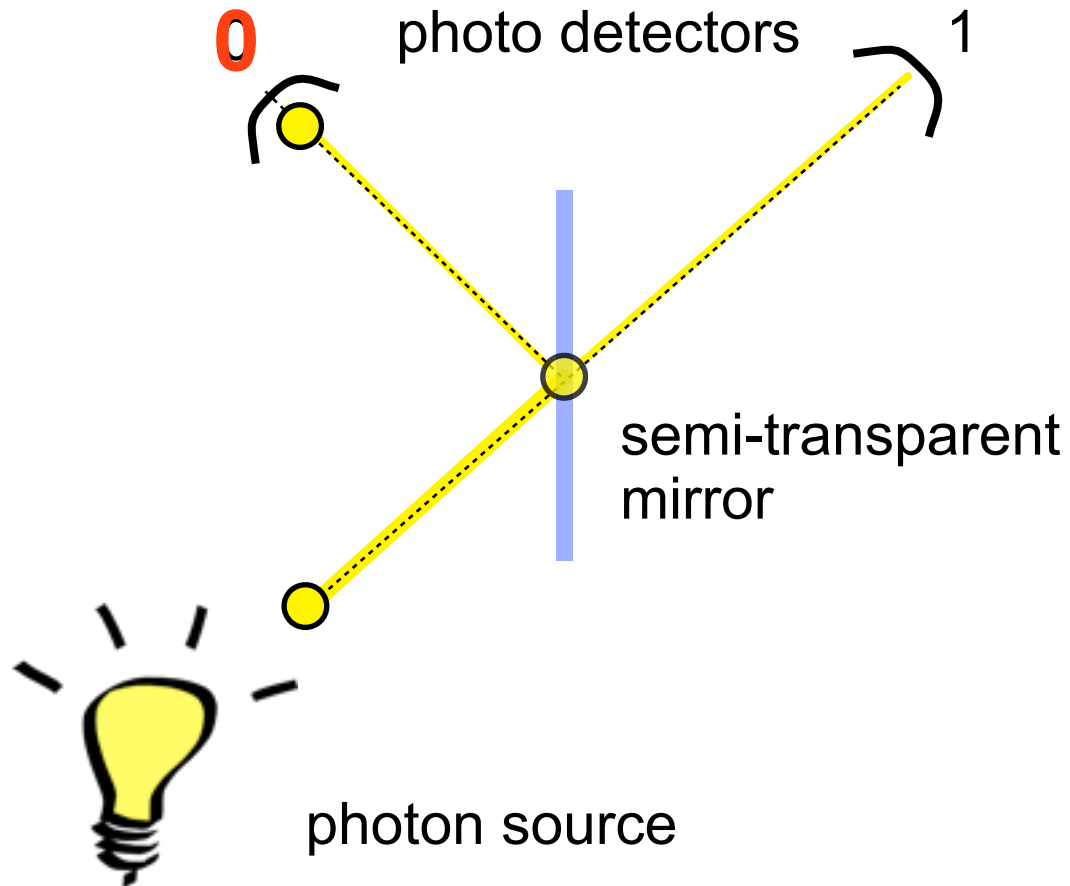
Quantum optimization(?)



Quantum computing



# True and perfect randomness



1. Photon source emits a photon
2. Photon hits semi-transparent mirror
3. Photon follows both paths
4. The photo detectors see the photon only in one place: **a random bit**

# The quantum bit (qubit)

Classical bits can be  $|0\rangle$  or  $|1\rangle$

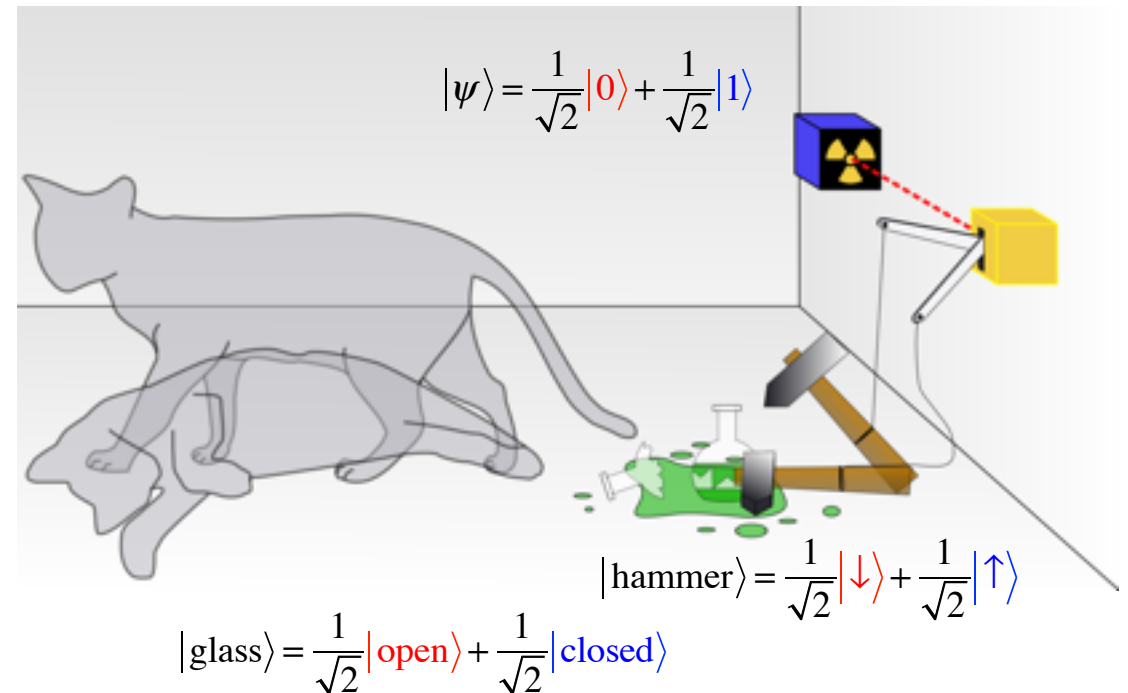
Qubits can be both at once

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$

“quantum superposition”

$$|\text{cat}\rangle = \frac{1}{\sqrt{2}}|\text{dead}\rangle + \frac{1}{\sqrt{2}}|\text{alive}\rangle$$

## Schrödinger's cat paradoxon



## Measuring a quantum superposition

- when measuring (looking) we only ever get one classical bit: 0 or 1

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle \quad \longrightarrow \quad \begin{array}{l} 0 \text{ with probability } |\alpha|^2 \\ 1 \text{ with probability } |\beta|^2 \end{array}$$

$|\alpha|^2 + |\beta|^2 = 1$

- When we look the cat is always either **dead** or **alive**!

- Quantum random number generator:

prepare and the state  $|\psi\rangle = \frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$  and measure

An application for a 1-qubit quantum computer!



# The incomprehensible magic of “quantum entanglement”

A single qubit gives a random bit when measured

$$|\psi\rangle = \frac{1}{\sqrt{2}}[|0\rangle + |1\rangle]$$

“Entangled states” can give random but identical results

$$|\psi\rangle = \frac{1}{\sqrt{2}}[|0\rangle_A |0\rangle_B + |1\rangle_A |1\rangle_B]$$

Measuring qubit  $A$  gives a random result  $a$

Measuring qubit  $B$  gives a random result  $b$

However, always  $a=b$  no matter how far apart the qubits are

A shared secret key that can be made provably secure!

## A serious restriction: no-cloning theorem

$$C|0\rangle \rightarrow |0\rangle|0\rangle$$

$$C|1\rangle \rightarrow |1\rangle|1\rangle$$

~~$$C|\psi\rangle \rightarrow C|\psi\rangle|\psi\rangle$$~~

A quantum state cannot be copied!

Bad news for quantum programmers

Excellent news for cryptographers



**NO CLONING!**

# Information content of a quantum register

A 2-qubit register

needs four complex numbers to be represented

but when measured only gives two bits of information

$$|\psi\rangle = \alpha_{00}|00\rangle + \alpha_{01}|01\rangle + \alpha_{10}|10\rangle + \alpha_{11}|11\rangle$$

An  $N$  qubit register

needs  $2^N$  complex numbers to be represented

but when measured only gives  $N$  bit of information

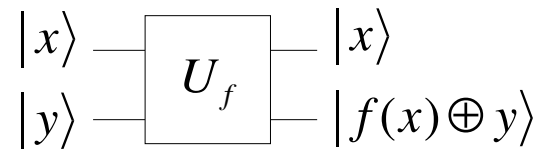
$$|\psi\rangle = \sum_{i_1, i_2, \dots, i_N} \alpha_{i_1 i_2 \dots i_N} |i_1 i_2 \dots i_N\rangle$$

Exponential intrinsic parallelism: operate on  $2^N$  inputs at once

But very limited readout of only  $N$  bits

# Calculating in superposition

Quantum computers can work on all possible inputs in superposition

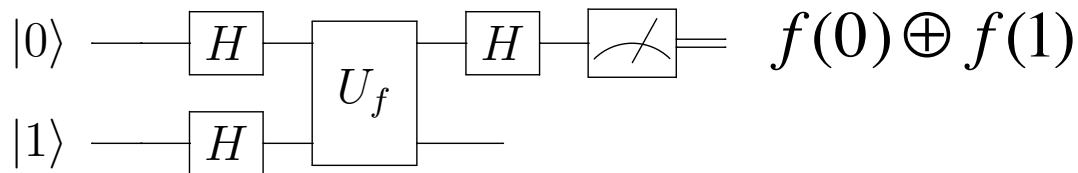


$$U_f |x\rangle |y\rangle \rightarrow |x\rangle |f(x) \oplus y\rangle$$

$$U_f (\alpha |0\rangle + \beta |1\rangle) |0\rangle \rightarrow \alpha |0\rangle |f(0)\rangle + \beta |1\rangle |f(1)\rangle$$

Measuring the result one only gets **either**  $f(0)$  **or**  $f(1)$ , chosen **randomly!**

Smartly compute one global result based on all inputs and measure it!



Determine whether  $f(0)=f(1)$   
with one function call

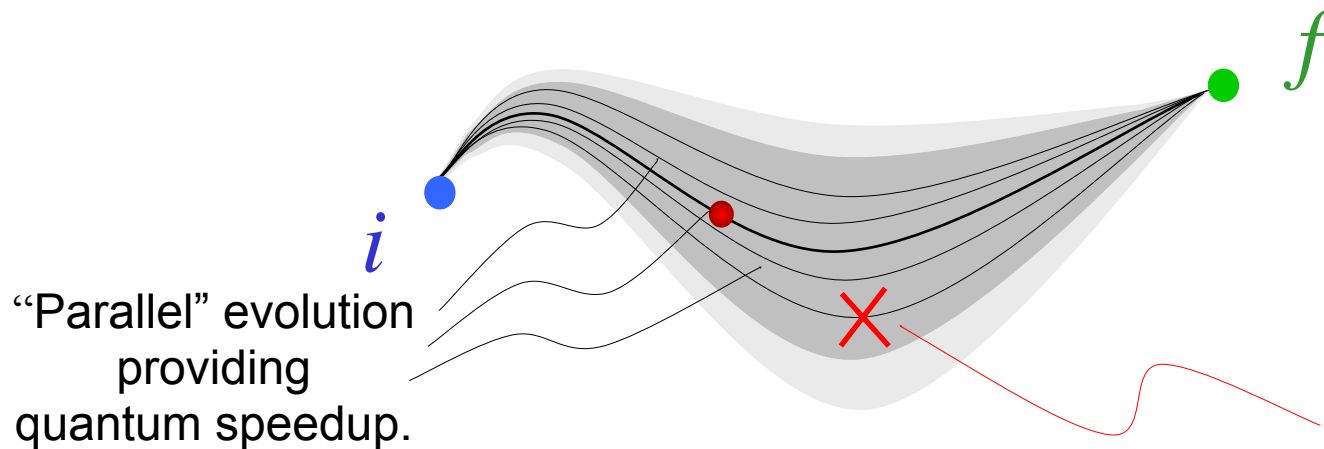
(Deutsch&Jozsa, 1992)

# Interlude: quantum hardware



# Observing the cat made it be either dead or alive!

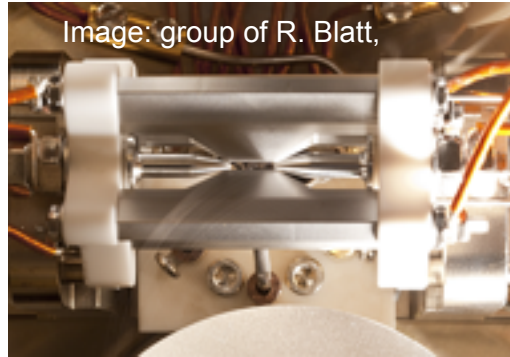
Qubits need to be well isolated from the environment!



Perturbations from the environment destroy the "parallel" quantum evolution of the computation

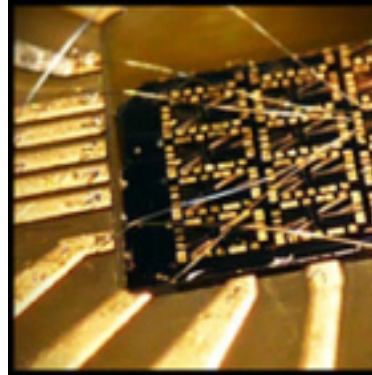


# Many different platforms

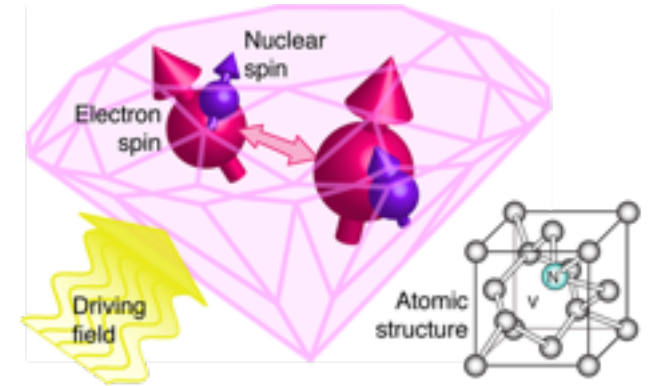


trapped ions  
20 qubits  
(R. Blatt, Innsbruck)

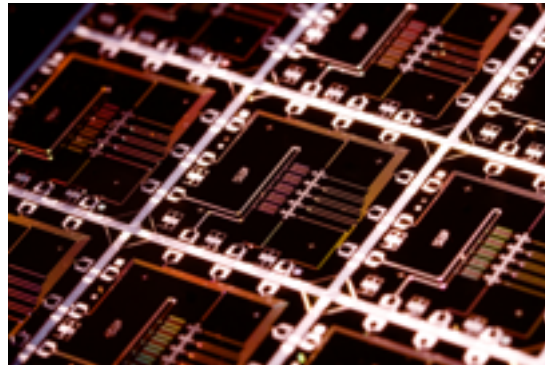
**100 gate operations on  
20 qubits**



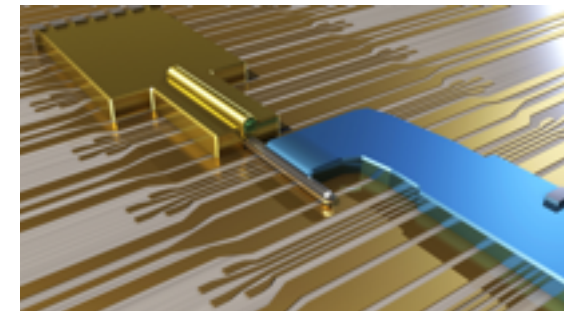
quantum dots  
(C. Marcus, Copenhagen)



defects in diamond



superconductors  
9 qubits  
(J. Martinis, UCSB)



Topological quantum bits  
(L. Kouwenhoven, Delft)

# Simulating quantum computers on classical computers

Simulating a quantum gate acting on  $N$  qubits needs  $O(2^N)$  memory and operations

Qubits	Memory	Time for one gate operation
10	16 kByte	microseconds on a watch
20	16 MByte	milliseconds on smartphone
30	16 GByte	seconds on laptop
40	16 TByte	minutes on supercomputer
50	16 PByte	hours on top supercomputer
60	16 EByte	long long time
80	size of visible universe	age of the universe

# Why should we build a quantum computer?

Simply because we can!

Somebody smart will figure out a use!

These arguments are not enough to justify the money it will cost

# Quantum computing beyond exa-scale

What are the important applications ...

... that we can solve on a quantum computer ...

... but not special purpose post-exa-scale classical hardware that we may build in ten years?



# What problems do we want to solve on a quantum computer?

design better drugs

cure cancer

counter climate  
change

fold proteins

fight hunger

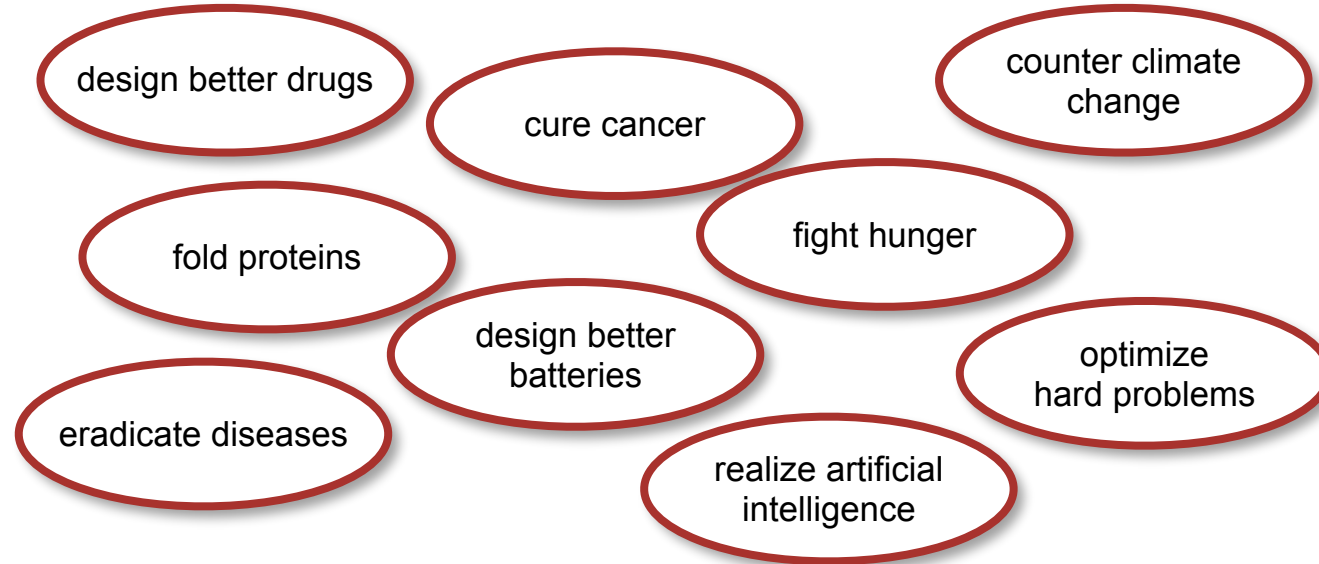
design better  
batteries

optimize  
hard problems

eradicate diseases

realize artificial  
intelligence

# What problems do we want to solve on a quantum computer?



This is a list for a quantum wishing well

Which of these can actually profit from quantum computers?



# A quantum machine to solve hard optimization problems

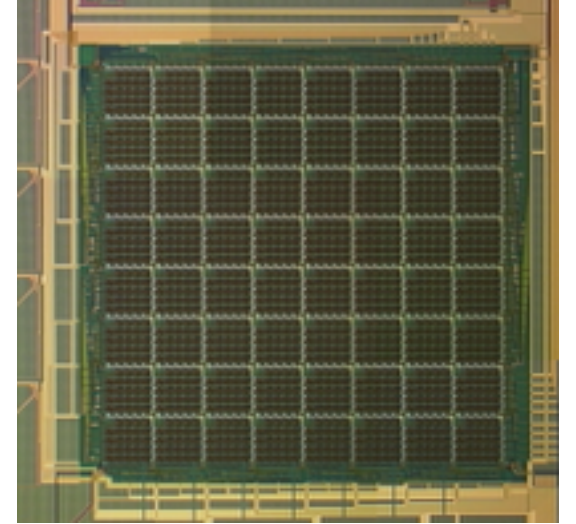


# The D-Wave quantum annealer

A device to solve quadratic binary optimization problems

$$C(x_1, \dots, x_N) = \sum_{ij} a_{ij} x_i x_j + \sum_i b_i x_i$$

with  $x_i = 0, 1$



Can be built with imperfect qubits

Significant engineering achievement to scale it to one thousand qubits

Nobody knows if it can solve NP-hard problems better than a classical computer

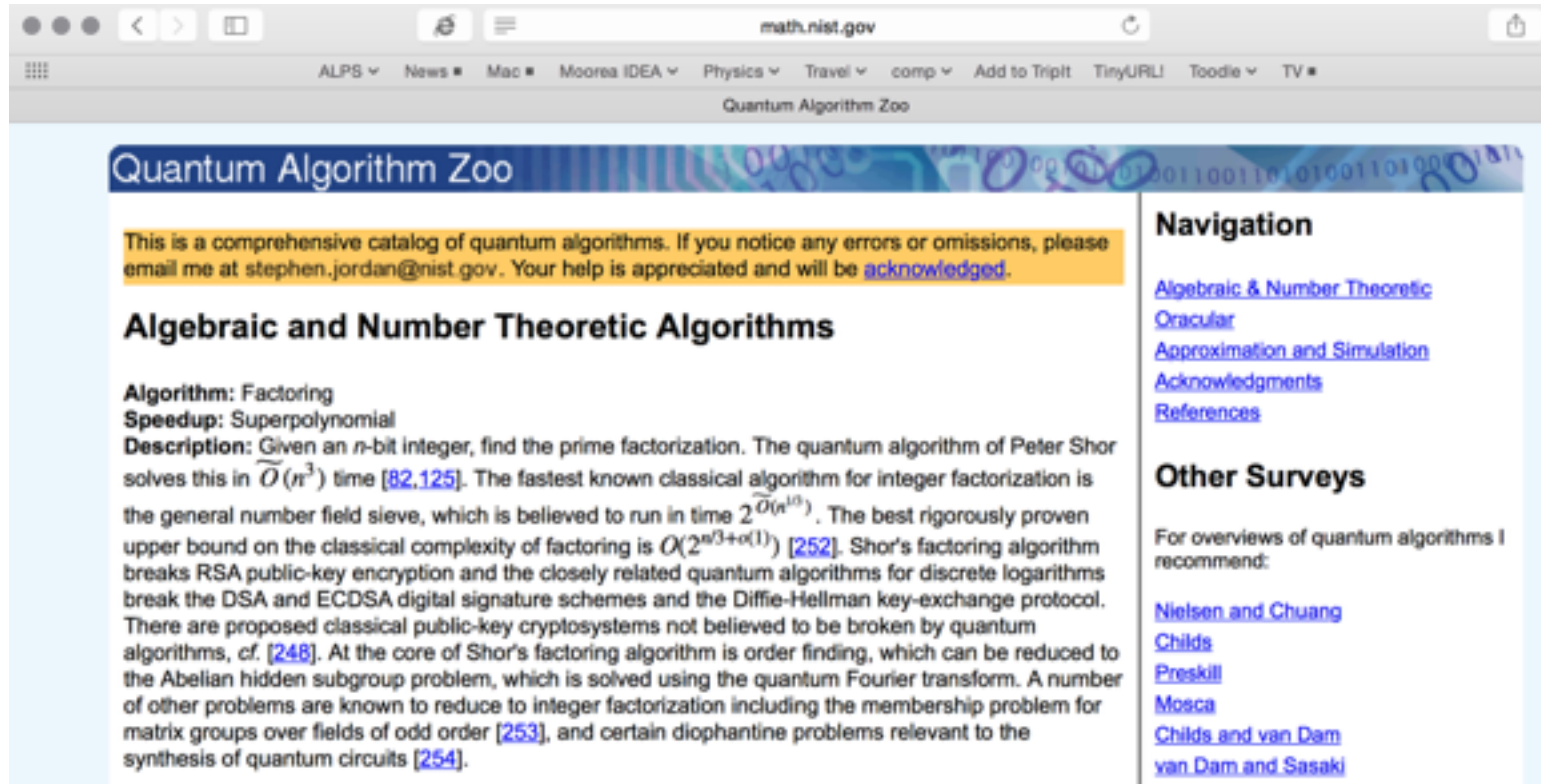
So far no scaling advantage has been observed

# Better look at algorithms with known quantum speedup

50+ quantum algorithms with known speedup

Can we use any of them in real-world applications?

<http://math.nist.gov/quantum/zoo/>



The screenshot shows a web browser window with the address bar displaying "math.nist.gov". The page title is "Quantum Algorithm Zoo". A navigation menu at the top includes "ALPS", "News", "Mac", "Moorea IDEA", "Physics", "Travel", "comp", "Add to Tripit", "TinyURL", "Toodler", and "TV". The main content area features a blue header with the text "Quantum Algorithm Zoo" and a yellow highlighted box containing the text: "This is a comprehensive catalog of quantum algorithms. If you notice any errors or omissions, please email me at [stephen.jordan@nist.gov](mailto:stephen.jordan@nist.gov). Your help is appreciated and will be [acknowledged](#)." Below this is the section "Algebraic and Number Theoretic Algorithms". Underneath, it lists "Algorithm: Factoring" and "Speedup: Superpolynomial". The "Description" text states: "Given an  $n$ -bit integer, find the prime factorization. The quantum algorithm of Peter Shor solves this in  $\tilde{O}(n^3)$  time [82,125]. The fastest known classical algorithm for integer factorization is the general number field sieve, which is believed to run in time  $2^{\tilde{O}(n^{1/3})}$ . The best rigorously proven upper bound on the classical complexity of factoring is  $O(2^{n^{0.75}})$  [252]. Shor's factoring algorithm breaks RSA public-key encryption and the closely related quantum algorithms for discrete logarithms break the DSA and ECDSA digital signature schemes and the Diffie-Hellman key-exchange protocol. There are proposed classical public-key cryptosystems not believed to be broken by quantum algorithms, cf. [248]. At the core of Shor's factoring algorithm is order finding, which can be reduced to the Abelian hidden subgroup problem, which is solved using the quantum Fourier transform. A number of other problems are known to reduce to integer factorization including the membership problem for matrix groups over fields of odd order [253], and certain diophantine problems relevant to the synthesis of quantum circuits [254]."

On the right side of the page, there is a "Navigation" section with links for "Algebraic & Number Theoretic", "Oracular", "Approximation and Simulation", "Acknowledgments", and "References". Below that is an "Other Surveys" section with the text "For overviews of quantum algorithms I recommend:" followed by links to "Nielsen and Chuang", "Childs", "Preskill", "Mosca", "Childs and van Dam", and "van Dam and Sasaki".

# Shor's algorithm for factoring

Factoring small numbers is easy:  $15 = 3 \times 5$

Factoring large numbers is hard classically:  $O(\exp(N^{1/3}))$  time for  $N$  digit-numbers

536939683642691194607950541533260051860418183893023116620231731884706135841697779  
 81247775554355964649044526158042091770292405381561410352725541976253778624830290  
 518096150501270434149272610204114236496946309670910771714302797950221151202416796  
 22849447805650987368350247829683054309216276674509735105639240298977591783205062  
 1619158848593319454766098482875128834780988979751083723214381986678381350567167

=

4363637625931498167701061252972058930130370651588109946621952523434903606572651613287  
 3421237667900245913537253744354928238018040554845306796065865605354860834270732796989  
 4210413710440109013191728001673

\*

1230486419064350262435007521990111788816176581586683476039159532309509792696707176253  
 0052007668467350605879541695798973080376300970096911310297914332946223591672260748684  
 8670728527914505738619291595079



Polynomial time on a quantum computer (P. Shor)

# Breaking RSA encryption with Shor's algorithm?

RSA	cracked in	CPU years	Shor
453 bits	1999	10	1 hour
768 bits	2009	2000	5 hours
1024 bits		1000000	10 hours



estimates based on 10 ns gate time and minimal number of  $2N+3$  qubits

Not a long-term “killer-app” since we can switch to post-quantum encryption

- quantum cryptography
- post-quantum encryption (e.g. lattice based cryptography)

# Grover search

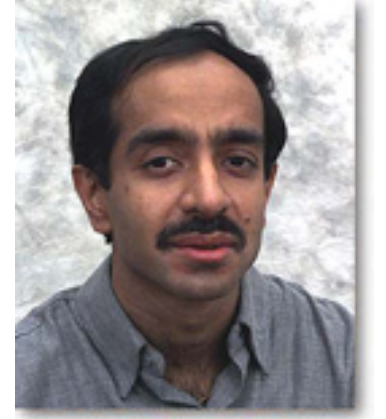
Search an unsorted database of  $N$  entries with  $\sqrt{N}$  queries

However, the query needs to be implemented!

- Querying an  $N$ -entry database needs at least  $O(N)$  hardware resources
- Can perform the query classically in  $\log(N)$  time given  $O(N)$  resources

Only useful if the query result can be efficiently calculated on the fly!

What are the important applications satisfying this criterion?





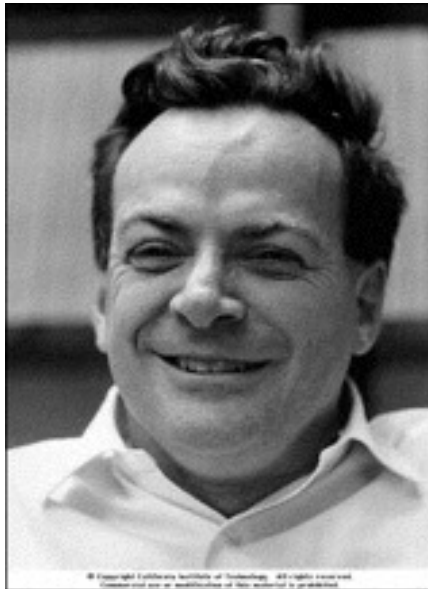
*International Journal of Theoretical Physics, Vol. 21, Nos. 6/7, 1982*

## Simulating Physics with Computers

**Richard P. Feynman**

*Department of Physics, California Institute of Technology, Pasadena, California 91107*

*Received May 7, 1981*



Feynman invented quantum computers to simulate quantum physics

We can surpass the best classical computers with only 50 qubits!

This will make physicists happy but is it enough to motivate



to to build one?

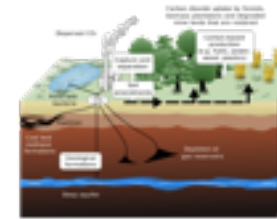
# First applications that reached a petaflop on Jaguar @ ORNL

Domain area	Code name	Institution	# of cores	Performance	Notes
Materials	DCA++	ORNL	213,120	<b>1.9 PF</b>	2008 Gordon Bell Prize Winner
Materials	WL-LSMS	ORNL/ETH	223,232	<b>1.8 PF</b>	2009 Gordon Bell Prize Winner
Chemistry	NWChem	PNNL/ORNL	224,196	1.4 PF	2008 Gordon Bell Prize Finalist
Materials	DRC	ETH/UTK	186,624	1.3 PF	2010 Gordon Bell Prize Hon. Mention
Nanoscience	OMEN	Duke	222,720	> 1 PF	2010 Gordon Bell Prize Finalist
Biomedical	MoBo	GaTech	196,608	780 TF	2010 Gordon Bell Prize Winner
Chemistry	MADNESS	UT/ORNL	140,000	550 TF	
Materials	LS3DF	LBL	147,456	442 TF	2008 Gordon Bell Prize Winner
Seismology	SPECFEM3D	USA (multiple)	149,784	165 TF	2008 Gordon Bell Prize Finalist

# Simulating quantum materials on a quantum computer

Can we use quantum computers to design new quantum materials?

- A room-temperature superconductor?
- Non-toxic designer pigments?
- A catalyst for carbon sequestration?
- Better catalysts for nitrogen fixation (fertilizer)?



Solving many materials challenges has

- exponentially complexity on classical hardware
- polynomial complexity on quantum hardware!

# Can quantum chemistry be performed on a small quantum computer?

Dave Wecker,<sup>1</sup> Bela Bauer,<sup>2</sup> Bryan K. Clark,<sup>2,3</sup> Matthew B. Hastings,<sup>2,1</sup> and Matthias Troyer<sup>4</sup>

<sup>1</sup>*Quantum Architectures and Computation Group, Microsoft Research, Redmond, WA 98052, USA*

<sup>2</sup>*Station Q, Microsoft Research, Santa Barbara, CA 93106-6105, USA*

<sup>3</sup>*Kavli Institute for Theoretical Physics, University of California, Santa Barbara, CA 93106, USA*

<sup>4</sup>*Theoretische Physik, ETH Zurich, 8093 Zurich, Switzerland*

Phys. Rev. A **90**, 022305 (2014)

Can a classically-intractable problem be solved  
on a small quantum computer?

Can a classically-intractable problem be solved on  
a huge quantum computer?

Can a classically-intractable problem be solved on  
the largest imaginable quantum computer?

# Simulating a quantum system on quantum computers

There are  $O(N^4)$  interaction terms in an  $N$ -electron system

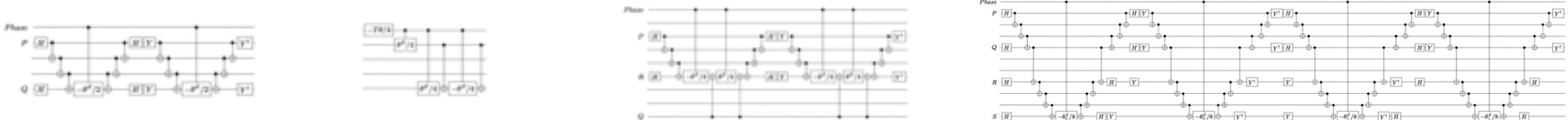
$$H = \sum_{pq} t_{pq} c_p^\dagger c_q + \sum_{pqrs} V_{pqrs} c_p^\dagger c_q^\dagger c_r c_s \equiv \sum_{m=1}^M H_m$$

$M = O(N^4)$  terms

We need to evolve separately under each of them

$$e^{-i\Delta t H} \approx \prod_{m=1}^M e^{-i\Delta \tau H_m}$$

Efficient circuits available for each of the  $N^4$  terms



Runtime estimates turn out to be  $O(NM^2) = O(N^9)$

# The polynomial time quantum shock

- Estimates for an example molecule: Fe<sub>2</sub>S<sub>2</sub> with 118 spin-orbitals

Gate count	10 <sup>18</sup>
Parallel circuit depth	10 <sup>17</sup>
Run time @ 10ns gate time	30 years

Quantum information theorists declare victory  
proving the existence of polynomial time algorithms

We need quantum software engineers to develop  
better algorithms and implementations



# The result of quantum software optimization

- Estimates for an example molecule: Fe<sub>2</sub>S<sub>2</sub> with 118 spin-orbitals

Gate count	10 <sup>18</sup>	Reduced gate count	10 <sup>11</sup>
Parallel circuit depth	10 <sup>17</sup>	Parallel circuit depth	10 <sup>10</sup>
Run time @ 10ns gate time	30 years	Run time @ 10ns gate time	2 minutes

- Attempting to reduce the horrendous runtime estimates we achieved  
*Wecker et al., PRA (2014), Hastings et al., QIC (2015), Poulin et al., QIC (2015)*
  - Reuse of computations:  $O(N)$  reduction in gates
  - Parallelization of terms:  $O(N)$  reduction in circuit depth
  - Optimizing circuits: 4x reduction in gates
  - Smart interleaving of terms: 10x reduction in time steps
  - Multi-resolution time evolution: 10x reduction in gates
  - Better phase estimation algorithms: 4x reduction in rotation gates

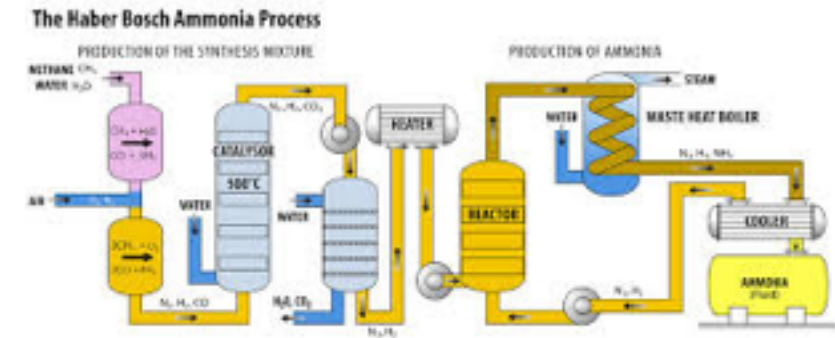
# Nitrogen fixation: a potential killer-app

Fertilizer production using Haber-Bosch process (1909)

Requires high pressures and temperatures

3-5% of the world's natural gas

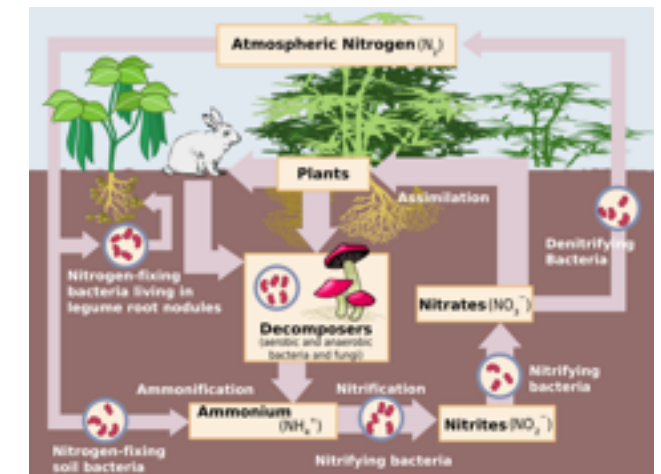
1-2% of the world's annual energy



But bacteria can do it cheaply at room temperature!

Quantum solution using about 400 qubits

- Understand how bacteria manage to turn air into ammonia
- Design a catalyst to enable inexpensive fertilizer production

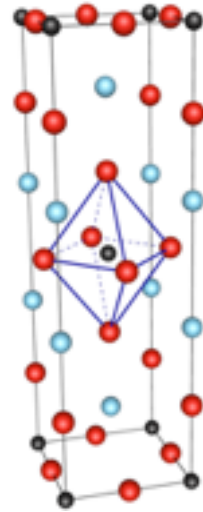


# What about a high temperature superconductor?

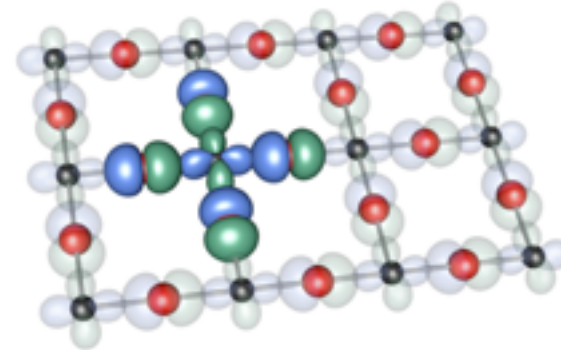
Orbitals per unit cell	$\approx 50$
Unit cells needed	$20 \times 20$
Number of orbitals	$N \approx 20'000$
Number of terms	$N^4$
Scaling of algorithm	$O(N^{5.5})$
Estimated runtime	age of the universe



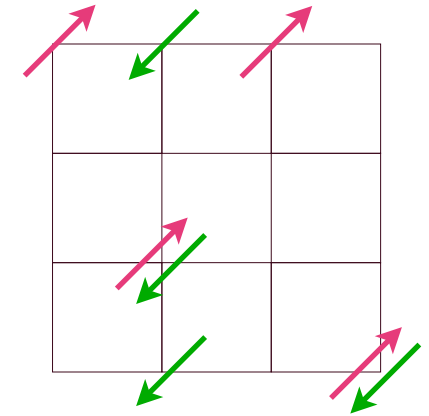
# Reduction to a simplified model



3D crystal structure



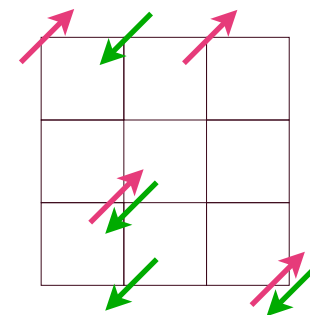
single 2D layer



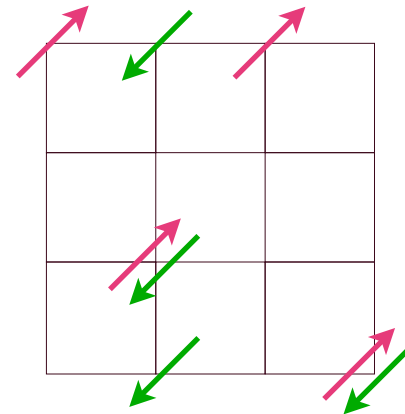
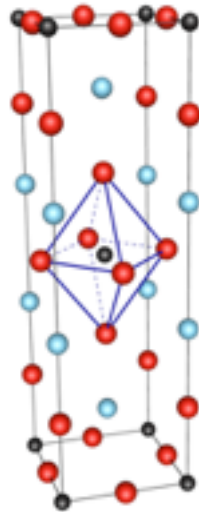
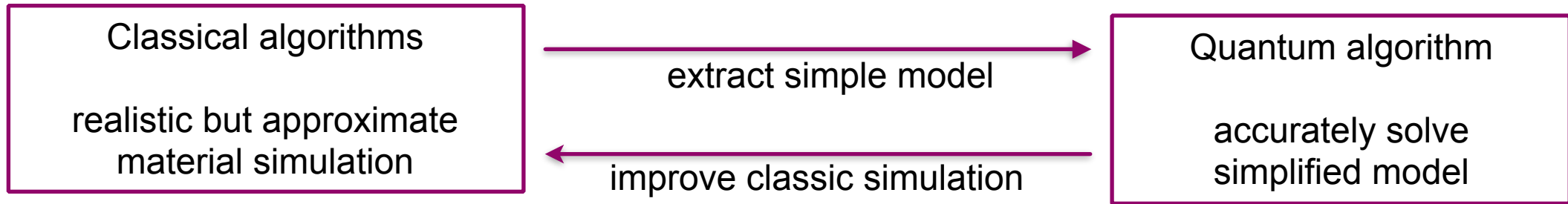
simplified model

# From materials to models on quantum computers

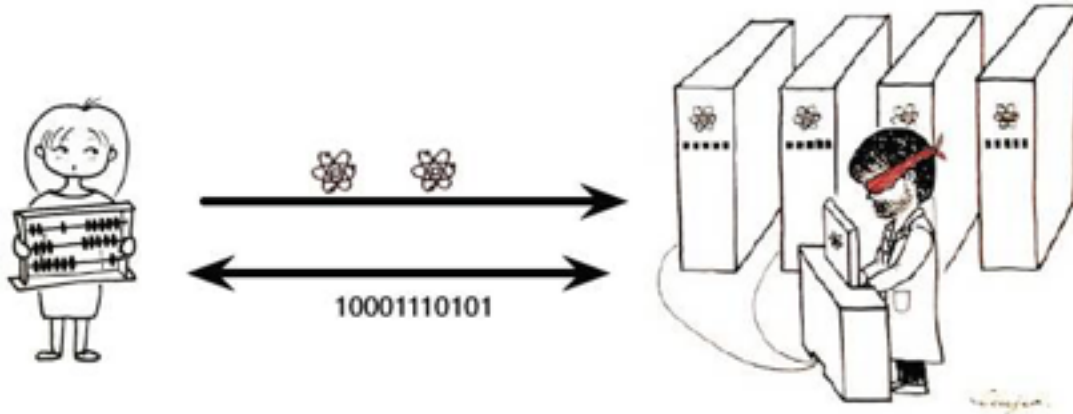
	Material	Model
Orbitals per unit cell	$\approx 50$	1
Unit cells needed	20x20	20x20
Number of orbitals	$N \approx 20'000$	$N \approx 800$
Number of terms	$N^4$	$O(N)$
Scaling of algorithm	$O(N^{5.5})$	$O(N^{0.5})$
Estimated runtime	age of the universe	seconds



# Hybrid quantum classical approaches



# There is much more!



## Blind quantum computing and search (Broadbent, Fitzsimons, Kashefi)

Cloud provides **cannot** know what the user does



## Quantum money (Aaronson, Farhi *et al*)



# What will we do with a quantum computer?

True random numbers with just one qubit

Secure communication with just a few qubits

Interesting real-world applications for a quantum computer

- Breaking of RSA encryption (?)
- Design of catalysts and materials
- Provably secure cloud computing



**We need quantum software engineers to explore more potential applications!**

# The quantum algorithms team



Dave Wecker



Matt Hastings

(Microsoft Research Redmond)



Nathan Wiebe



Bela Bauer

(Microsoft Station Q Santa Barbara)



Chetan Nayak



David Poulin  
(Sherbrooke)



Andrew Doherty  
(Sydney)



Bryan Clark  
(UIUC)



Andy Millis  
(Columbia)