

Influence Maximization: Pushing the Limits of Combinatorial Optimizations and Online Learning

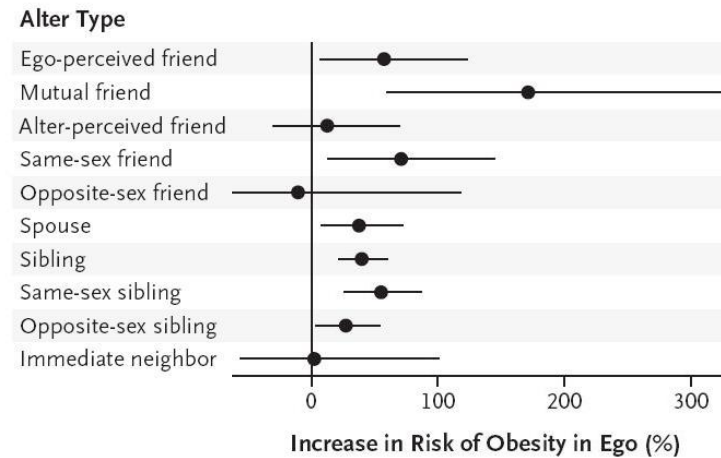
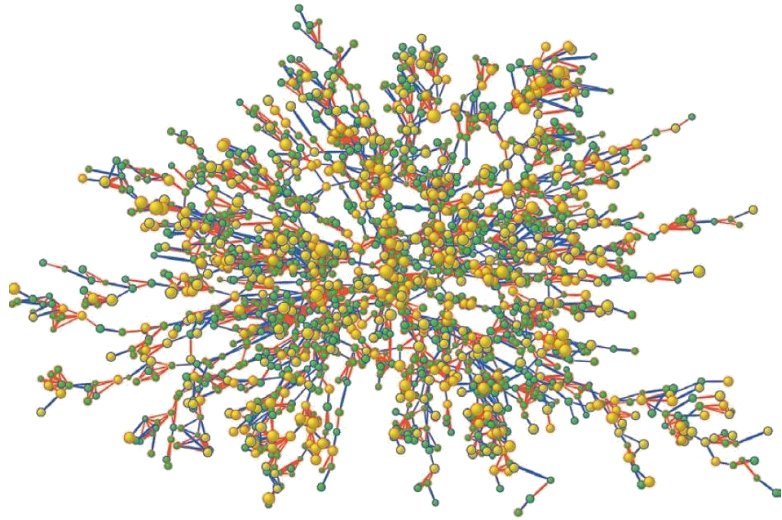
Wei Chen 陈卫

Microsoft Research Asia

Information and Influence Propagation in Networks



Examples of Studies on Influence in Networks: Obesity and Stopping Smoking



Christakis N A and Fowler J H. The spread of obesity in a large social network over 32 years. *New England Journal of Medicine*, 2007(357.4):370~379

Christakis N A and Fowler J H. The collective dynamics of smoking in a large social network. *New England Journal of Medicine*, 2008(358.21):2249~2258

Voting Mobilization: A Facebook Study

- Voting mobilization [Bond et al, Nature'2012]
 - show a facebook msg. on voting day with faces of friends who voted
 - generate 340K additional votes due to this message, among 60M people tested

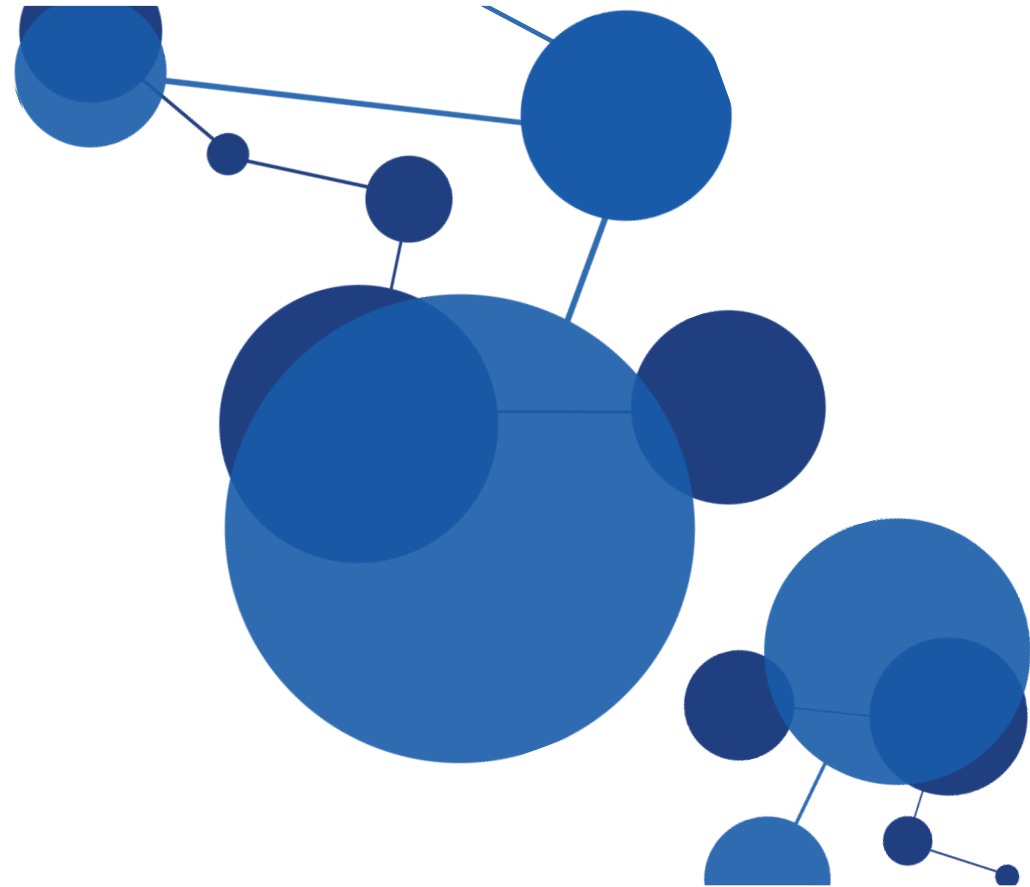


Bond R M, Fariss C J, Jones J J, Kramer A D I, Marlow C, Settle J E, and Fowler J H. A 61-million-person experiment in social influence and political mobilization. *Nature*, 2012(489):295~298

Influence Propagation Modeling and Optimizations

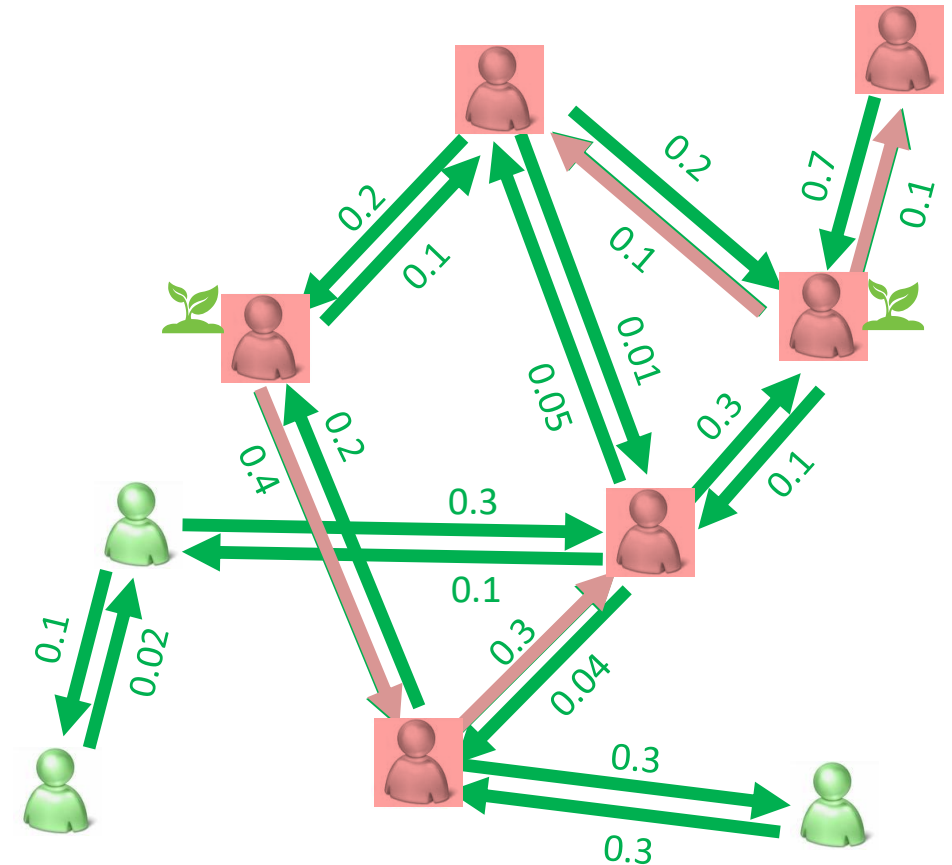
- How to model influence propagation in a social network?
 - Stochastic diffusion models
- How to optimize the influence propagation effect?
 - Influence maximization and its variants
- One core problem: **Influence maximization**
 - Find a small number of individuals in a network to generate a large influence
 - Applications in viral marketing, diffusion monitoring, rumor control, etc.

Influence Maximization in a Nutshell



Independent Cascade (IC) Model

- Each edge (u, v) has an *influence probability* $p(u, v)$
- Initially seed nodes in S_0 are activated
- At each step t , each node u activated at step $t - 1$ activates its neighbor v independently with probability $p(u, v)$
- **Influence spread $\sigma(S)$** : expected number of activated nodes
- Other models: linear threshold (LT), triggering, general threshold, etc.



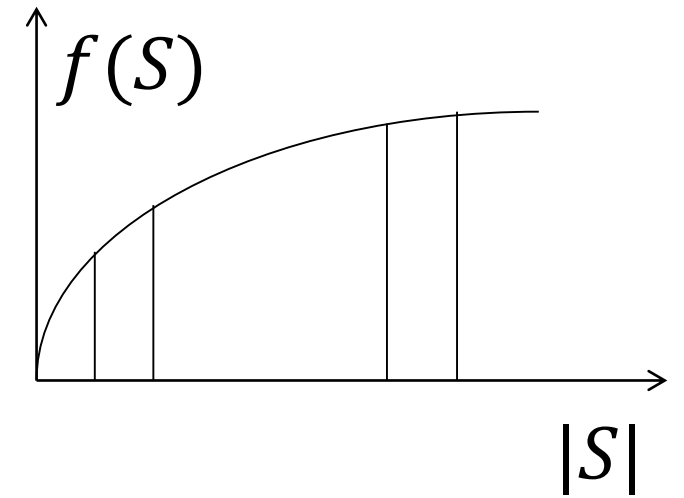
Influence Maximization

- Given a social network, a diffusion model with given parameters, and a number k , find a seed set S of at most k nodes such that the influence spread of S is maximized.
- Based on *submodular function* maximization
- [Kempe, Kleinberg, and Tardos, KDD'2003]

Kempe D, Kleinberg J M, and Tardos É. Maximizing the spread of influence through a social network. KDD'2003

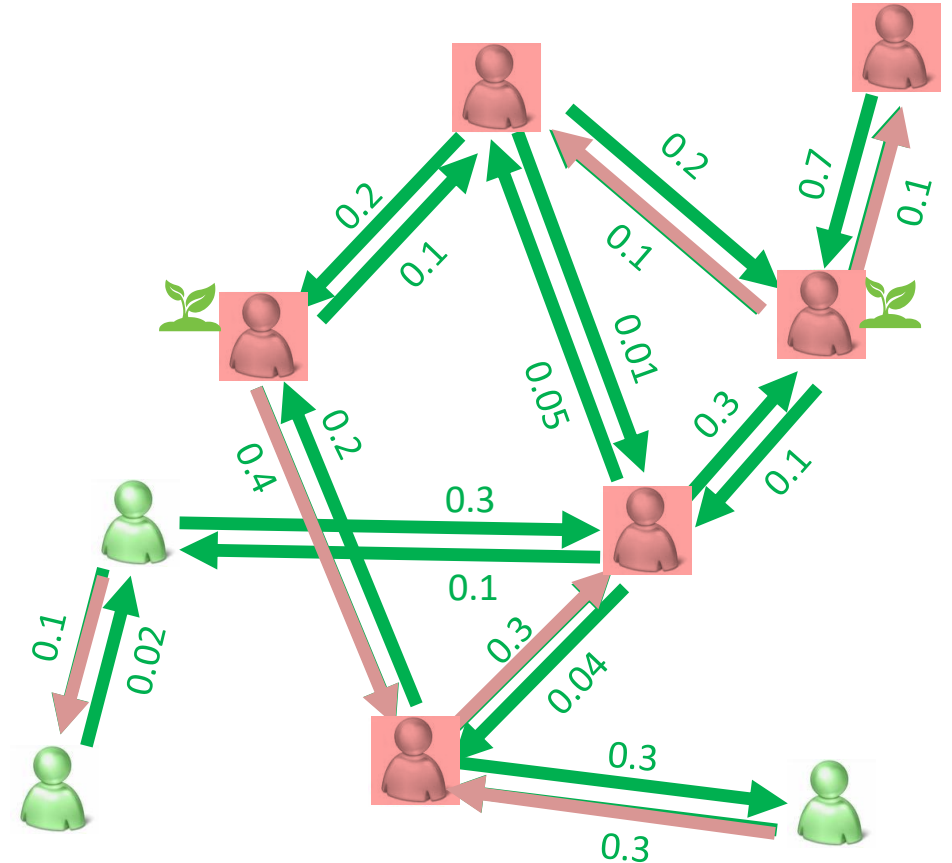
Submodular Set Functions

- **Sumodularity** of set functions $f: 2^V \rightarrow R$
 - for all $S \subseteq T \subseteq V$, all $v \in V \setminus T$,
$$f(S \cup \{v\}) - f(S) \geq f(T \cup \{v\}) - f(T)$$
 - diminishing marginal return
 - an equivalent form: for all $S, T \subseteq V$
$$f(S \cup T) + f(S \cap T) \leq f(S) + f(T)$$
- **Monotonicity** of set functions f : for all $S \subseteq T \subseteq V$,
$$f(S) \leq f(T)$$



Submodularity of Influence Spread Function $\sigma(S)$

- Independent cascade model is equivalent to
 - sample **live edges** by edge probabilities
 - activate nodes reachable from S in the **live-edge graph**
- $\sigma(S) = \sum_L \Pr\{L\} \cdot |\Gamma(L, S)|$
 - $\Gamma(L, S)$: set of nodes reachable from S in live-edge graph L
 - $|\Gamma(L, S)|$ is a coverage function, easy to show it is submodular



Greedy Algorithm for Submodular Function Maximization

- 1: initialize $S = \emptyset$;
- 2: for $i = 1$ to k do
- 3: select $u = \operatorname{argmax}_{w \in V \setminus S} [f(S \cup \{w\}) - f(S)]$
- 4: $S = S \cup \{u\}$
- 5: end for
- 6: output S

Property of the Greedy Algorithm

- Theorem: If the set function f is monotone and submodular with $f(\emptyset) \geq 0$, then the greedy algorithm achieves $(1 - 1/e)$ approximation ratio, that is, the solution S found by the greedy algorithm satisfies:

$$f(S) \geq \left(1 - \frac{1}{e}\right) \max_{S' \subseteq V, |S'|=k} f(S')$$

- [Nemhauser, Wolsey and Fisher, 1978]

Nemhauser G L, Wolsey L A, and Fisher M L. An analysis of approximations for maximizing submodular set functions. Mathematical Programming 1978

Challenges and Research Coverage

- Scalability challenge:
 - In IC (and LT) models, computing influence spread $\sigma(\mathcal{S})$ for any given \mathcal{S} is #P-hard [Chen et al. KDD'2010, ICDM'2010], and Monte Carlo simulation is slow
 - Scalable influence maximization
- Adaptivity challenge:
 - Can we adapt to partial feedbacks? --- adaptive influence maximization
- Learning challenge:
 - How to learn the diffusion model?
 - How to use online feedback for optimization --- online influence maximization
- Complex model challenge:
 - Other variants of influence diffusion models --- competitive and complementary influence maximization, non-submodular influence maximization, etc.

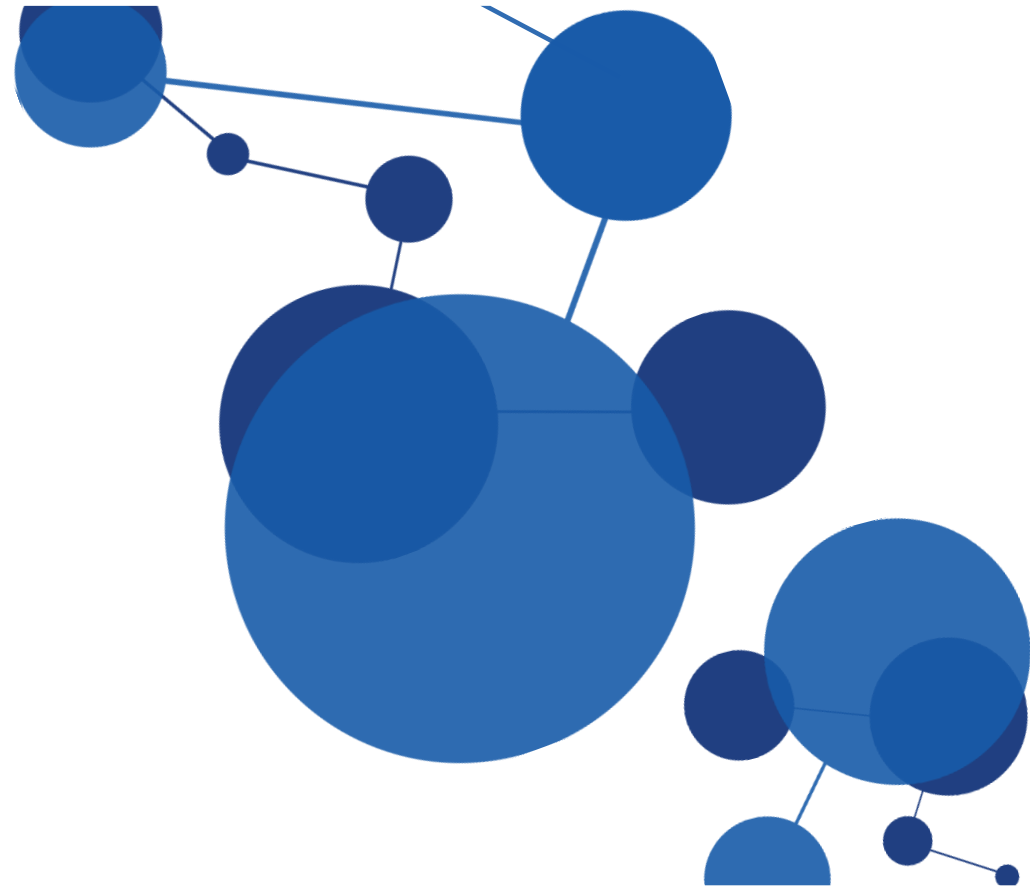
Pushing the Limits of Optimization and Learning

- Influence maximization
 - Sitting at the boundary of feasibility
- Examples discussed in this talk
 - Adaptive influence maximization: new variants in adaptive maximization

	Adaptive submodular	Non-adaptive submodular
Independent feedback	prior studies	IC + myopic feedback
Dependent feedback	IC + full-adoption feedback	LT + myopic/full-adoption feedback

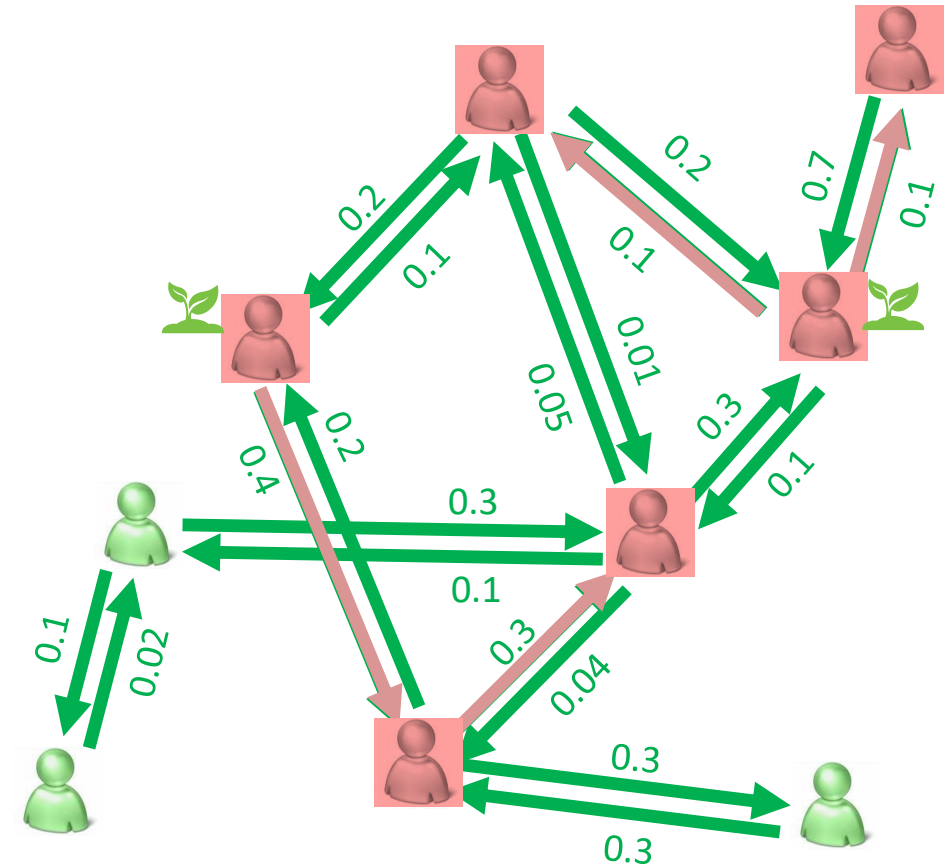
- Online influence maximization --- general combinatorial multi-armed bandit framework
 - prior studies: linear rewards, exact offline oracle
 - Online IM: nonlinear rewards, approximation oracle, probabilistically triggered arms

Adaptive Influence Maximization



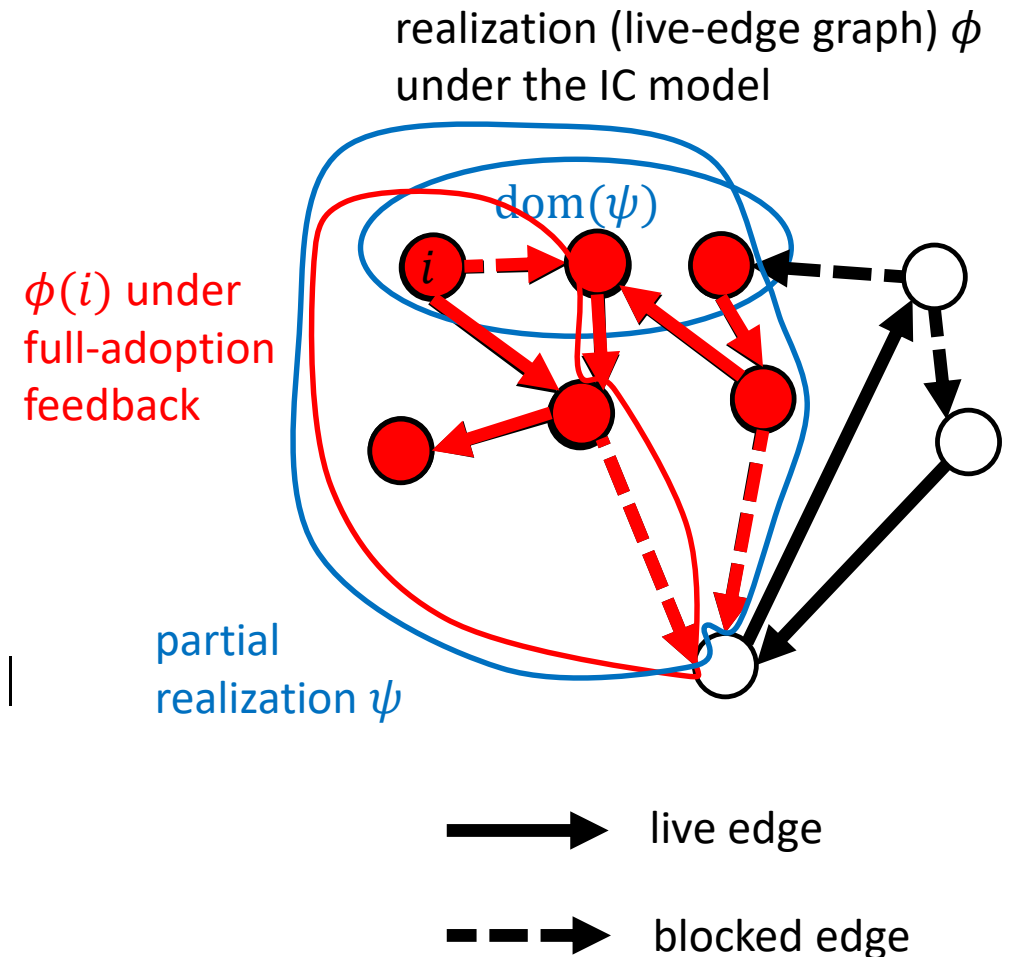
Adaptive Influence Maximization: Model

- Influence propagation: IC Model
- Seed selection: one-by-one, instead of a batch of k nodes
 - After selecting each seed node, obtain feedback on the propagation from the seed --- can be used to help subsequent seed selection
- Two feedback models
 - Full-adoption feedback: all downstream propagation from the selected seed, whether an edge passes through influence, whether a node is activated
 - Myopic feedback: only immediate propagation from the seed to its out-neighbors are given as the feedback



Adaptive Submodularity

- Realization ϕ : all randomness in a propagation (random live-edge graph)
- Partial realization ψ : feedback collected (partial propagation) from the currently selected seeds $\text{dom}(\psi)$
- Adaptive Submodularity: a node u 's marginal influence is higher on a smaller partial realization than on a larger partial realization
 - If $\psi \subseteq \psi'$, $\Delta(u|\psi') \leq \Delta(u|\psi)$
- Adaptive Monotonicity: a node u 's marginal influence on any partial realization is nonnegative
 - $\Delta(u|\psi) \geq 0$, as long as ψ has non-zero probability to occur



Important Results on Adaptive Submodularity

- **Approximation ratio**
 - Adaptive Greedy Algorithm
 - On every step, greedily select the next entry with the largest marginal influence: select $v = \operatorname{argmax}_u \Delta(u|\psi)$
 - If the model is adaptive monotone and adaptive submodular, adaptive greedy algorithm is a $1 - 1/e$ approximation of the adaptive optimal solution. [GK11]
- **Adaptivity gap**: supremum ratio of the adaptive optimal vs. non-adaptive optimal: $\sup_{G,k} \frac{OPT_A(G,k)}{OPT_N(G,k)}$
 - If the model is adaptive monotone and adaptive submodular, and the feedback are mutually independent, the adaptivity gap is at most $\frac{e}{e-1}$. [AN16]

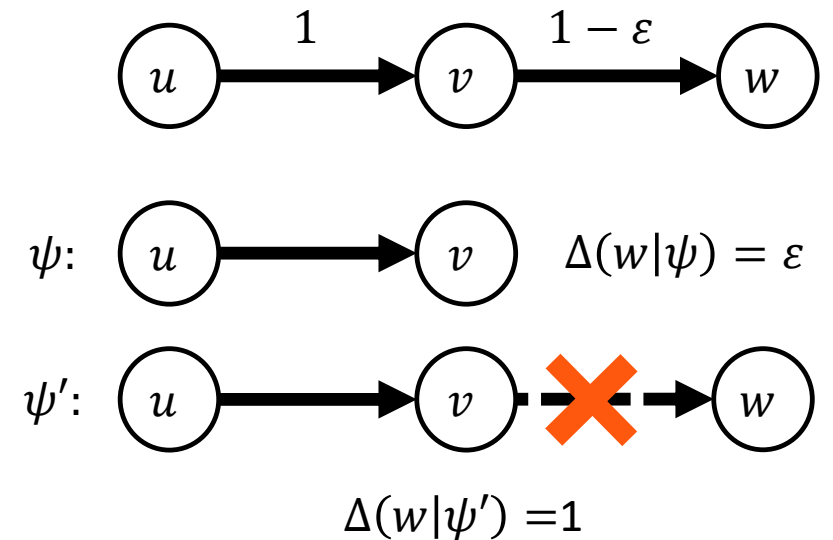
Golovin D and Krause A. Adaptive submodularity: theory and applications in active learning and stochastic optimization. Journal of Artificial Intelligence Research, 2011

Asadpour A and Nazerzadeh H. Maximizing stochastic monotone submodular functions. Management Science, 2016

Adaptive Submodularity on Influence Maximization

	IC model	LT model
Full-adoption feedback	adaptive submodular	not adaptive submodular
Myopic feedback	not adaptive submodular	not adaptive submodular

- IC+myopic not adaptive submodular
 - Example to the right
 - [GK11] conjectures that adaptive greedy is still a constant approximation
 - We answer this conjecture affirmatively [NeurIPS'19]



Adaptivity Gap on Influence Maximization

- Related to adaptive Submodularity and feedback independence

	Adaptive submodular	Non-adaptive submodular
Independent feedback	$\frac{e}{e-1}$ [AN16]	IC + myopic feedback $\left[\frac{e}{e-1}, 4\right]$ [PC19]
Dependent feedback	IC + full-adoption feedback Partial answers on specific graphs [CP19]: In-arborescences: $\left[\frac{e}{e-1}, \frac{2e}{e-1}\right]$ Out-arborescences: $\left[\frac{e}{e-1}, 2\right]$ Bipartite graphs: $\frac{e}{e-1}$	LT + myopic/full-adoption feedback ? Triggering + full-adoption: unbounded [CPST20] Triggering+myopic ?

Peng B and [Chen W](#). Adaptive influence maximization with myopic feedback, NeurIPS'2019

[Chen W](#) and Peng B. On adaptivity gaps of influence maximization under the independent cascade model with full adoption feedback. ISAAC'2019

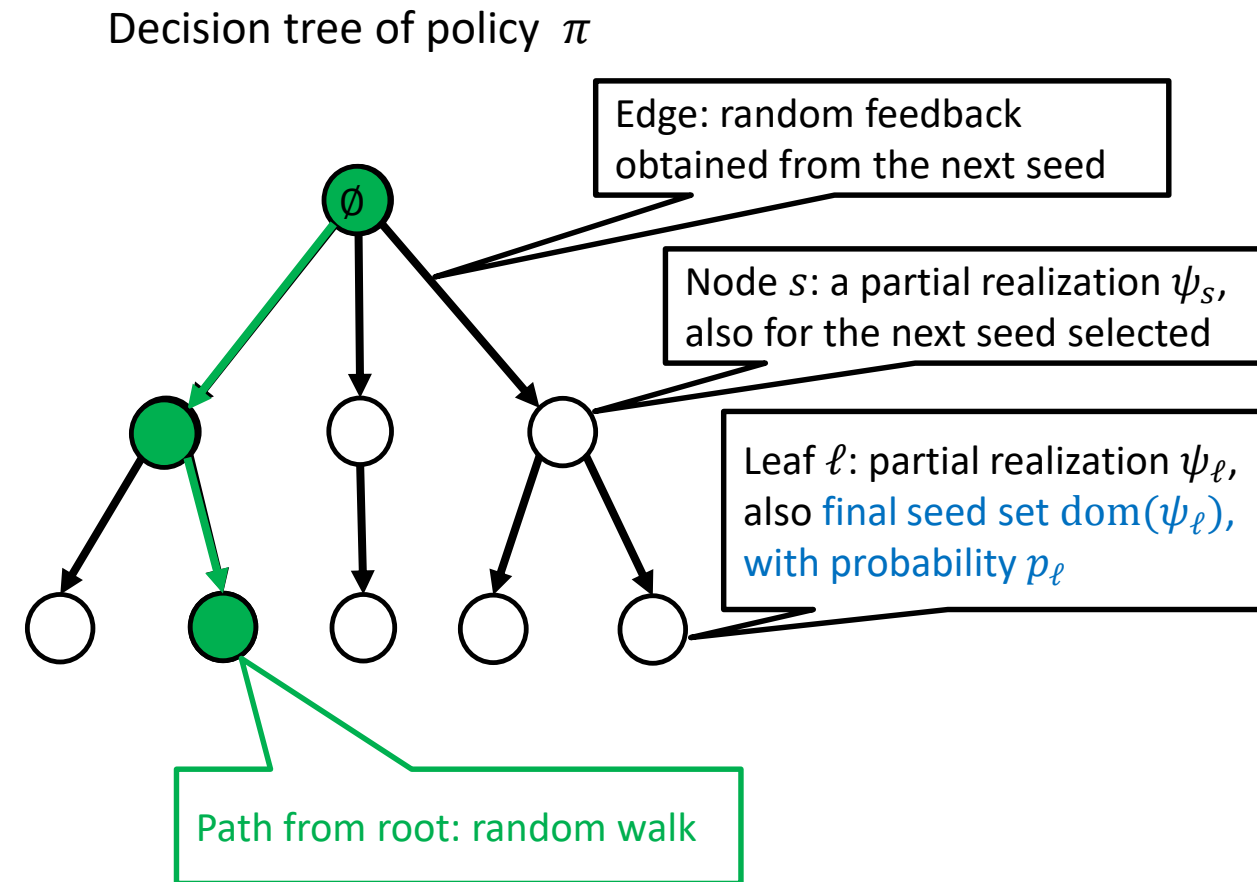
[Chen W](#), Peng B, Schoenebeck G, and Tao B. Adaptive greedy versus non-adaptive greedy for influence maximization. AAAI'2020

Implications on IC + Myopic Feedback

- Adaptive greedy is $\frac{1}{4} \cdot \left(1 - \frac{1}{e}\right)$ approximation of the adaptive optimal solution
 - Adaptive greedy is $\left(1 - \frac{1}{e}\right)$ approximation of the non-adaptive optimal
 - Non-adaptive optimal is $\frac{1}{4}$ of the adaptive optimal
- Answers the open conjecture by Golovin and Krause (2011)
- First study on adaptive maximization with a non-adaptive submodular model

Idea on the Analysis, IC + Myopic, $\text{Gap} \leq 4$

- Decision tree representation of adaptive policy π
- Random walk non-adaptive policy $\mathcal{W}(\pi)$: Select $\text{dom}(\psi_\ell)$ with probability p_ℓ
- Fictitious hybrid policy $\bar{\pi}$ and aggregate adaptive influence spread $\bar{\sigma}(\bar{\pi})$
- Show: $\sigma(\pi) \leq \bar{\sigma}(\bar{\pi}) \leq 4\sigma(\mathcal{W}(\pi))$



Fictitious Hybrid Policy $\bar{\pi}$ and Aggregate Adaptive Influence Spread $\bar{\sigma}(\bar{\pi})$

- Work simultaneously on three independent realizations Φ_1, Φ_2, Φ_3
- $\bar{\pi}$ selects seeds adaptively exactly like π working on Φ_1
- But for each selected seed u , it has three independent chances to activate its out-neighbors, according to Φ_1, Φ_2, Φ_3
- The expected number of activated nodes is the aggregate adaptive influence spread $\bar{\sigma}(\bar{\pi})$
- Obviously, $\sigma(\pi) \leq \bar{\sigma}(\bar{\pi})$

Connecting Aggregate Adaptive Spread $\bar{\sigma}(\bar{\pi})$ with Non-Adaptive Spread $\sigma(\mathcal{W}(\pi))$

- t -th aggregate influence spread $\sigma^t(\mathcal{S})$ and t -th aggregate adaptive influence spread $\sigma^t(\pi)$, $t = 1, 2, 3$
 - each seed gets t independent chances of activating its out-neighbors
 - $\bar{\sigma}(\bar{\pi}) = \sigma^3(\pi)$, $\sigma(\mathcal{W}(\pi)) = \sigma^1(\mathcal{W}(\pi))$
- Represent $\sigma^t(\mathcal{W}(\pi))$ and $\sigma^t(\pi)$ by non-adaptive marginal gains $\Delta_{ft}(u|\text{dom}(\psi_s))$ and adaptive marginal gains $\Delta_{ft}(u|\psi_s)$, respectively
 - telescoping series on node s along a path in the decision tree
- $\Delta_{f^3}(u|\psi_s) \leq 2\Delta_{f^2}(u|\text{dom}(\psi_s)) \Rightarrow \sigma^3(\pi) \leq 2\sigma^2(\mathcal{W}(\pi))$
 - key lemma, crucially relying on (a) feedback independence and (b) (non-adaptive) submodularity of influence utility function on live-edge graphs
- $\sigma^2(\mathcal{W}(\pi)) \leq 2\sigma(\mathcal{W}(\pi))$

Greedy Adaptivity Gap

- Motivation:
 - optimal solutions cannot be reached
 - Practical algorithms are mostly greedy-based
- Greedy Adaptivity Gap: ratio between adaptive greedy vs. non-adaptive greedy
- Results:

	IC model	LT model	Triggering model (more general)
Full-adoption	tight lower bound: $1 - 1/e$	tight lower bound: $1 - 1/e$	Upper bound: unbounded
Myopic	tight lower bound: $1 - 1/e$	tight lower bound: $1 - 1/e$	

Chen W, Peng B, Schoenebeck G, and Tao B. Adaptive greedy versus non-adaptive greedy for influence maximization. AAI'2020

Many Open Problems

- Adaptivity gap:

	IC model	LT model	Triggering model (more general)
Full-adoption	result on general graphs? tighter result for special graphs?	?	unbounded
Myopic	$\left[\frac{e}{e-1}, 4\right]$, tight result?	?	?

- Greedy adaptivity gap:

	IC model	LT model	Triggering model (more general)
Full-adoption	upper bound?	upper bound?	lower: $1 - 1/e$ upper: unbounded
Myopic	upper bound $\leq \frac{4e}{e-1}$ tight upper bound?	upper bound?	upper bound?

- Better adaptive algorithms than greedy?

Online Influence Maximization

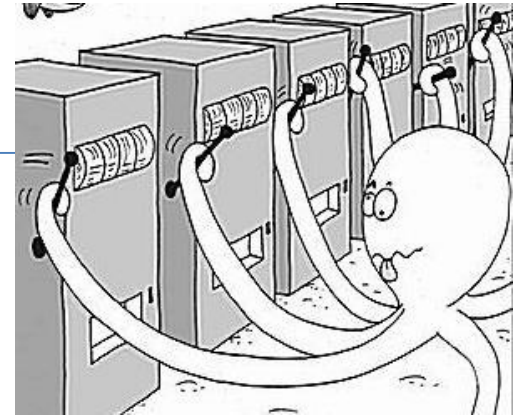


Online Influence Maximization

- Edge influence probabilities are unknown, need to be learned
- Multiple rounds of online influence maximization. In each round,
 - select k seeds to influence the network
 - observe the diffusion paths and results
 - collect the reward --- the number of nodes activated
 - use the observed feedback to update learning statistics, which is used for seed selection in later rounds
- Falls into the online learning (multi-armed bandit) framework

Multi-Armed Bandit Problem

- There are m arms (machines)
- Arm i has an unknown reward distribution on $[0,1]$ with unknown mean μ_i
 - best arm $\mu^* = \max \mu_i$
- In each round, the player selects one arm to play and observes the reward
- Performance metric: Regret:
 - Regret after playing T rounds $= T\mu^* - \mathbb{E}[\sum_{t=1}^T R_t(i_t^A)]$
- Objective: minimize regret in T rounds
- Balancing exploration-exploitation tradeoff
 - exploration: try new arms
 - exploitation: keep playing the best arm so far
- Wide applications: Any scenario requiring selecting best choice from online feedback
 - online recommendations, advertising, wireless channel selection, social networks, A/B testing



Classical MAB Algorithm: UCB1

- 1: for each arm i : $\hat{\mu}_i = 1$ (empirical mean), $T_i = 0$ (number of observation)
 - 2: for $t = 1, 2, 3, \dots$ do
 - 3: for each arm i : $\rho_i = \sqrt{\frac{3 \ln t}{2T_i}}$ (confidence radius)
 - 4: for each arm i : $\bar{\mu}_i = \min\{\hat{\mu}_i + \rho_i, 1\}$ (upper confidence bound, UCB)
 - 5: $j = \operatorname{argmax}_i \bar{\mu}_i$
 - 6: play arm j , observe its reward $X_{j,t}$
 - 7: update $\hat{\mu}_j = (\hat{\mu}_j \cdot T_j + X_{j,t}) / (T_j + 1)$; $T_j = T_j + 1$
 - 6: end-for
-
- For exploration
- For exploitation

Guarantee of the UCB1 Algorithm

- Finite-horizon regret:
 - distribution dependent: $O\left(\sum_{\Delta_i > 0} \frac{1}{\Delta_i} \ln T\right)$, $\Delta_i = \mu^* - \mu_i$
 - distribution independent: $O(\sqrt{mT \ln T})$
- [Auer, Cesa-Bianchi, and Fischer, 2002]

Auer P, Cesa-Bianchi N, and Fischer P. Finite-time analysis of the multiarmed bandit problem. *Machine Learning Journal*, 2002(47.2-3):235~256

Challenges Applying UCB1 to Online IM

- exponential number of seed sets
 - cannot treat each seed set as an arm
- non-linear reward functions
- offline problem is already NP-hard
- probabilistically triggering new arms in a play

Extending the MAB Framework

- Extend MAB to combinatorial MAB framework with probabilistically triggered arms (CMAB-T)
 - Model: In each round one **action/super-arm** is played, which triggers a set of **base arms** (triggering may be probabilistic)
 - **precisely characterize the bounded smoothness condition required to solve CMAB-T**
 - propose the CUCB algorithm based on an offline approximation oracle
 - **distribution-dependent and distribution-independent regret analysis**
 - applicable to a large class of combinatorial online learning problems
- **[Chen et al JMLR'2016, Wang and Chen, NIPS'2017]**

Chen W, Wang Y, Yuan Y, and Wang Q. Combinatorial multi-armed bandit and its extension to probabilistically triggered arms. Journal of Machine Learning Research, 2016

Wang Q and **Chen W**. Improving regret bounds for combinatorial semi-bandits with probabilistically triggered arms and its applications. NIPS'2017

CUCB Algorithm

- 1: for each arm i : $\hat{\mu}_i = 1$ (empirical mean), $T_i = 0$ (number of observation)
- 2: for $t = 1, 2, 3, \dots$ do
- 3: for each arm i : $\rho_i = \sqrt{\frac{3 \ln t}{2T_i}}$ (confidence radius)
- 4: for each arm i : $\bar{\mu}_i = \min\{\hat{\mu}_i + \rho_i, 1\}$ (upper confidence bound, UCB)
- 5: $\mathcal{S} = \text{OfflineOracle}(\bar{\mu}_1, \dots, \bar{\mu}_m)$
- 6: play action/super-arm \mathcal{S} , observe triggered arm outcomes $\{X_{j,t}\}$
- 7: for each observed j : update $\hat{\mu}_j = (\hat{\mu}_j \cdot T_j + X_{j,t}) / (T_j + 1)$; $T_j = T_j + 1$
- 6: end-for

Regret Bounds

- $O\left(\sum_i \frac{1}{\Delta_{\min}^i} B_1^2 K \ln T\right)$ distribution-dependent regret
 - i : base arm index
 - B_1 : one-norm bounded-smoothness constant
 - K : maximum number of arms any action can trigger
 - T : time horizon, total number of rounds
 - Δ_{\min}^i : minimum gap between α fraction of the optimal reward and the reward of any action that could trigger arm i (α is the offline approximation ratio)
- $O(B_1 \sqrt{mKT \ln T})$ distribution-independent regret
- For influence maximization, B_1 is the largest number of nodes any node can reach
- Main technical contribution: (a) characterizing the exact conditions (b) more sophisticated techniques in the analysis

Open Problems in Online Influence Maximization

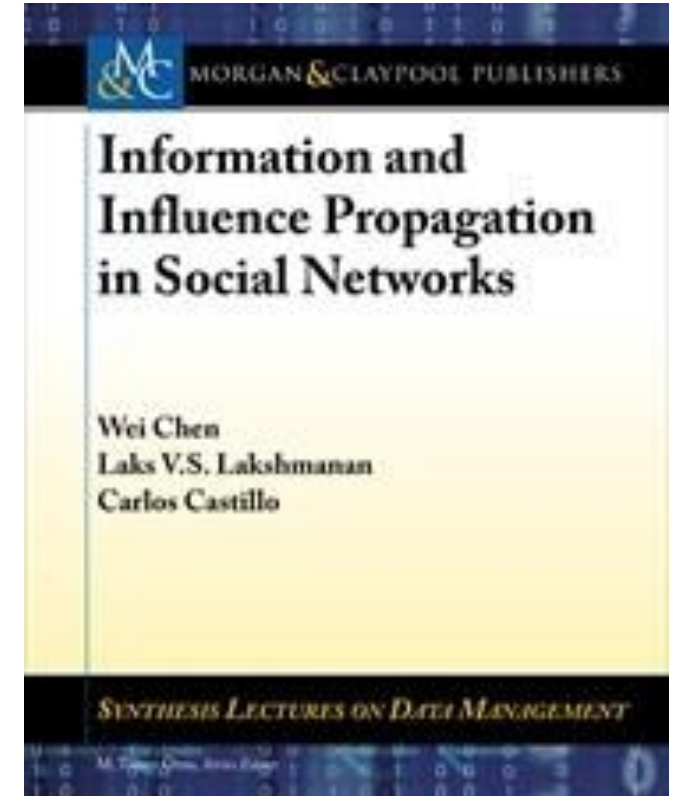
- More efficient algorithms specific to influence maximization?
- Regret lower bound?
- Online IM with nonstationary distribution?

Conclusion and Future Work

- Influence maximization is a rich ground for studying many optimization and learning tasks
 - Right at the boundary of feasibility --- pushing the tasks to new limits
 - Other directions beyond adaptive and online influence maximization:
 - Scalable algorithms (well studied)
 - Learnability of influence functions (some studies)
 - Optimization from samples (initial study)
 - Non-submodular influence maximization (some studies)
 - Influence maximization + game theory (some initial studies)

Reference Resources

- Search “Wei Chen Microsoft”
 - Monograph: “Information and Influence Propagation in Social Networks”, Morgan & Claypool, 2013
 - my papers and talk slides
 - My upcoming book (in Chinese): 大数据网络传播模型和算法 (Network Diffusion Models and Algorithms for Big Data)



Thanks!

