DATA PARALLEL PROGRAMMING IN HASKELL An Overview

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Ubiquitous Parallelism

It's parallelism

...but not as we know it!

Our goals

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Exploit parallelism of commodity hardware easily:

- ▶ Performance is important, but...
- ‣ …productivity is more important.

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Exploit parallelism of commodity hardware easily:

- ▶ Performance is important, but...
- ‣ …productivity is more important.
- Semi-automatic parallelism
	- ‣ Programmer supplies a parallel algorithm
	- ‣ No explicit concurrency (no concurrency control, no races, no deadlocks)

Computer Vision [Haskell 2011] **Edge detection**

Canny on 2xQuad-core 2.0GHz Intel Harpertown

Canny (512x512)

Physical Simulation [Haskell 2012a] Fluid flow

Runtimes for Jos Stam's Fluid Flow Solver

We can beat C!

Fluid flow (1024x1024)

GPU beats the CPU (includes all transfer times)

Medical Imaging [Haskell 2012a]

Interpolation of a slice though a 256 × 256 × 109 × 16-bit data volume of an MRI image

Functional Parallelism

Our ingredients

- Control effects, not concurrency
- Types guide data representation and behaviour
- Bulk-parallel aggregate operations

Our ingredients

Haskell is by default pure

Control effects, not concurrency

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Purity & Types

Purity and parallelism

 processList :: [Int] -> ([Int], Int) processList list = (sort list, maximum list)

Purity and parallelism **processList :: [Int] -> ([Int], Int) processList list = (sort list, maximum list)** function argument function body

*** Purity: function result depends only on arguments**

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Parallelism: execution order only constrained by explicit data dependencies

λ By default pure ϵ

Types track purity

λ By default pure (2)

 $Pure = no effects$ Impure = may have effects

λ By default pure (2)

λ By default pure (2)

λ By default pure \odot

Types guide execution

For bulk-parallel, aggregate operations, we introduce a new datatype:

Array r sh e

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- \triangleright D delayed array (represented as a function)
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- Shape: dimensionality of the array
	- ‣ DIM0, DIM1, DIM2, and so on
- Element type: stored in the array
	- ‣ Primitive types (Int, Float, etc.) and tuples

zipWith :: (Shape sh, Source r1 a, Source r2 b)
\n
$$
\Rightarrow
$$
 (a -> b -> c)
\n \Rightarrow Array r1 sh a
\n \Rightarrow Array r2 sh b
\n \Rightarrow Array D sh c

dotp $v w = sumAll (zipWith (*) w w)$

type Vector r e = Array r DIM1 e

 dotp :: (Num e, Source r1 e, Source r2 e) => Vector r1 e -> Vector r2 e -> e

dotp $v = sumAll (zipWith (*) v w)$

Elements are any type of numbers…

dotp $v = sumAll (zipWith (*) v w)$ type Vector r e = Array r DIM1 e dotp :: (Num e, Source r1 e, Source r2 e) => Vector r1 e -> Vector r2 e -> e

Data Parallelism

Parallelism

ABSTRACTION

Parallelism is safe for pure functions (i.e., functions without external effects)

Concurrency

Data Parallelism

Collective operations have got a single conceptual thread of control

Parallelism

ABSTRACTION

Parallelism is safe for pure functions (i.e., functions without external effects)

Concurrency

Types & Embedded Computations

Core i7 970 CPU

NVIDIA GF100 GPU

COARSE-GRAINED VERSUS FINE-GRAINED PARALLELISM

GPUs require careful program tuning

Core i7 970 CPU

NVIDIA GF100 GPU

✴**SIMD: groups of threads executing in lock step (warps)**

✴**Need to be careful about control flow**

COARSE-GRAINED VERSUS FINE-GRAINED PARALLELISM

GPUs require careful program tuning

24,576 THREADS

Core i7 970 CPU

NVIDIA GF100 GPU

✴**SIMD: groups of threads executing in lock step (warps)**

✴**Need to be careful about control flow**

✴**Latency hiding: optimised for regular memory access patterns**

✴**Optimise memory access**

COARSE-GRAINED VERSUS FINE-GRAINED PARALLELISM

GPUs require careful program tuning

Code generation for embedded code

Embedded DSL

- ‣ Restricted control flow
- ‣ First-order GPU code

Generative approach based on combinator templates

Code generation for embedded code

Embedded DSL

- **Exercited control flow ✓ limited control structures**
- **Eirst-order GPU code**
- Generative approach based on combinator templates

Code generation for embedded code

Embedded DSL

- **Exercited control flow ✓ limited control structures**
- ▶ First-order GPU cod

Generative approach based on combinator templates **✓ hand-tuned access patterns**

[DAMP 2011]

Dot product

```
dotp :: Vector Float -> Vector Float 
     -> Acc (Scalar Float)
dotp xs ys 
  = let
      xs' = use xs ys' = use ys
     in
    fold (+) 0 (zipWith (*) xs' ys')
```


Nested Data Parallelism

Modularity

- Standard (Fortran, CUDA, etc.) is flat, regular parallelism
- Same for our libraries for functional data parallelism for multicore CPUs (Repa) and GPUs (Accelerate)
- But we want more…

smvm :: SparseMatrix -> Vector -> Vector smvm sm $v = [$: sumP (dotp sv v) | sv <- sm :]

Nested parallelism

*** Modular**

* Irregular, nested data structures

- ‣ Sparse structures, tree structures
- ‣ Hierachrchical decomposition
- Nesting to arbitrary, dynamic depth: divide & conquer
- Lots of compiler work: still very experimental! [FSTTCS 2008, ICFP 2012, Haskell 2012b]

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- Types identify restricted code for specialised hardware, such as GPUs
- Types guide parallelising program transformations

Summary

Blog: <http://justtesting.org>[/](http://justtesting.org) Twitter: @TacticalGrace

- Core ingredients
	- ‣ Control purity, not concurrency
	- **Fypes guide representations and behaviours**
	- ‣ Bulk-parallel operations
- Get it
	- ‣ Latest Glasgow Haskell Compiler (GHC)
	- ‣ Repa, Accelerate & DPH packages from Hackage (Haskell library repository)

Thank you!

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